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MADALGO seminar by Zhewei Wei, Aarhus University

Space Complexity of 2-Dimensional Approximate Range Counting

Abstract:

We study the problem of 2-dimensional orthogonal range counting with absolute additive error. Given a set P of n points drawn from an $n \times n$ grid and an error parameter ϵ , the goal is to build a data structure, such that for any orthogonal range R , the data structure can return the number of points in $P \cap R$ with additive error $\epsilon \cdot n$. A well-known solution for this problem is the ϵ -approximation. Informally speaking, an ϵ -approximation of P is a subset A of P that allows us to estimate the number of points in $P \cap R$ by counting the number of points in $A \cap R$. It is known that an ϵ -approximation of size $O(1/\epsilon \log^{2.5} 1/\epsilon)$ exists for any P with respect to orthogonal ranges, and the best lower bound is $\Omega(1/\epsilon \log 1/\epsilon)$.

The ϵ -approximation is a rather restricted data structure, as we are not allowed to store any information other than the coordinates of a subset of points in P . In this talk, we explore what can be achieved without any restriction on the data structure. We first describe a data structure that uses $O(1/\epsilon \log 1/\epsilon \log \log 1/\epsilon \log n)$ bits that answers queries with error $\epsilon \cdot n$. We then prove a lower bound that any data structure that answers queries with error $O(\log n)$ must use $\Omega(n \log n)$ bits. This lower bound has two consequences: 1) answering queries with error $O(\log n)$ is as hard as answering the queries exactly; and 2) our upper bound cannot be improved in general by more than an $O(\log \log 1/\epsilon)$ factor.

Joint work with Ke Yi